

# Quantum GANs A Physics Perspective

Roberto Cappuccio

Dipartimento di Scienze Fisiche, della Terra e dell'Ambiente

Università di Siena

# Index

- Classical Gan
- Classical GAN Sample Code
- Quantum GAN
- Quantum GAN Sample Code
- Quantum GAN in Physics

# What's a GAN

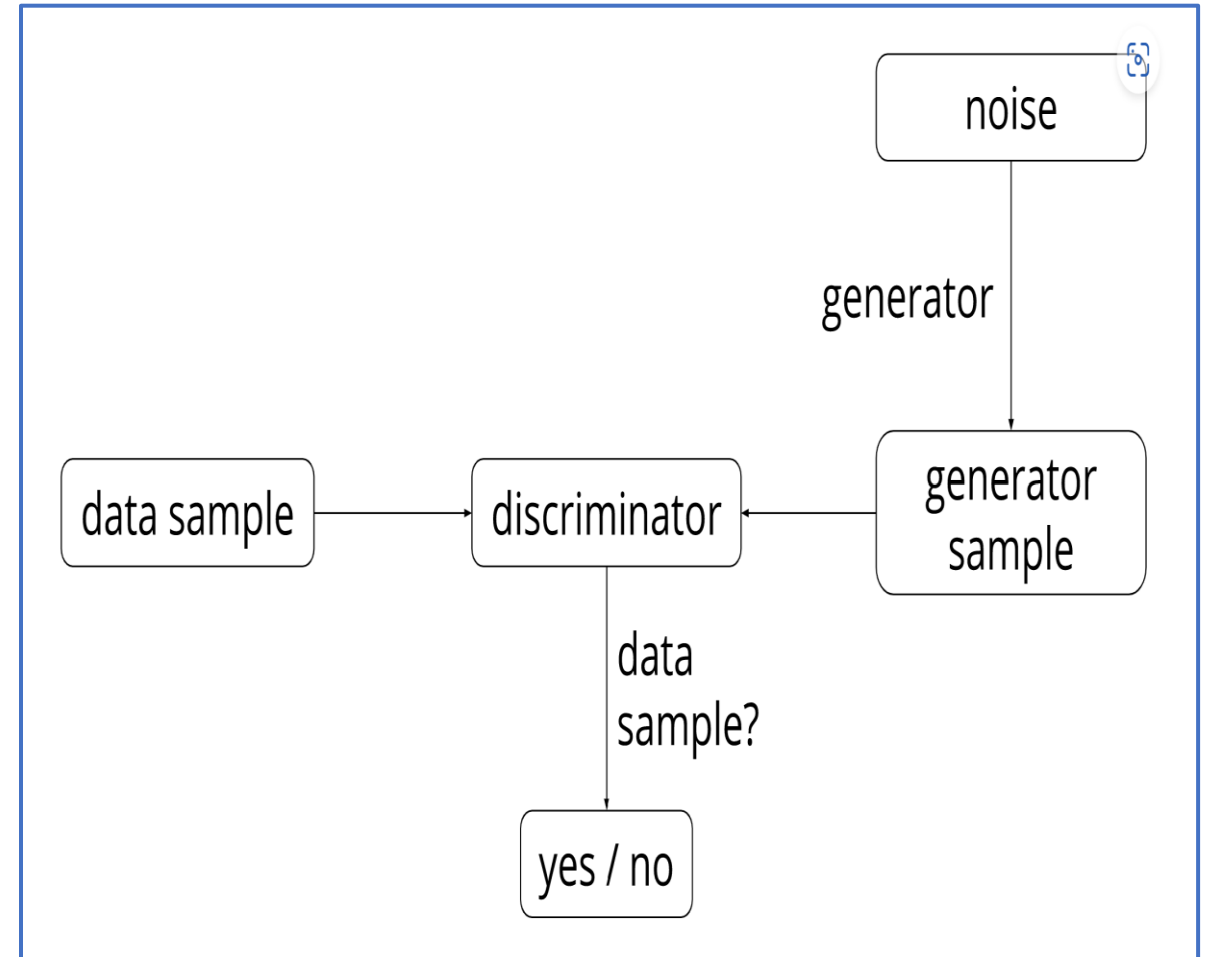
## GAN stands for Generative Adversarial Network

- The specific goal of GANs is to generate data from scratch
- Typical domain is Images but not only (Videos, Music, Texts, Functions, Simulation Data etc)
- GANs are made up by 2 main blocks:
  - **GENERATOR** (Leo): it generates fake data from noise distribution, by pretending they are true
  - **DISCRIMINATOR** (Tom): it wants to prove the generated data are fake



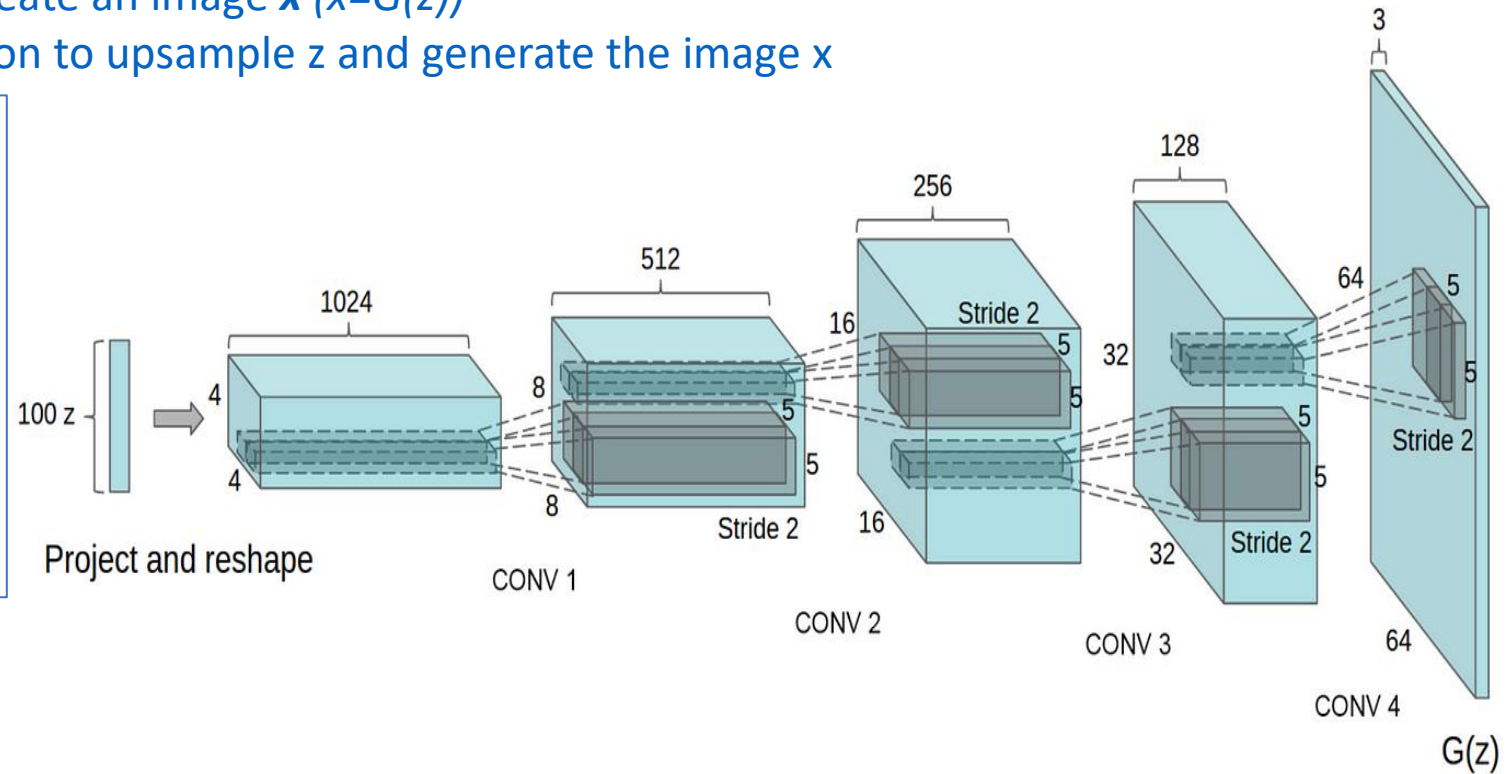
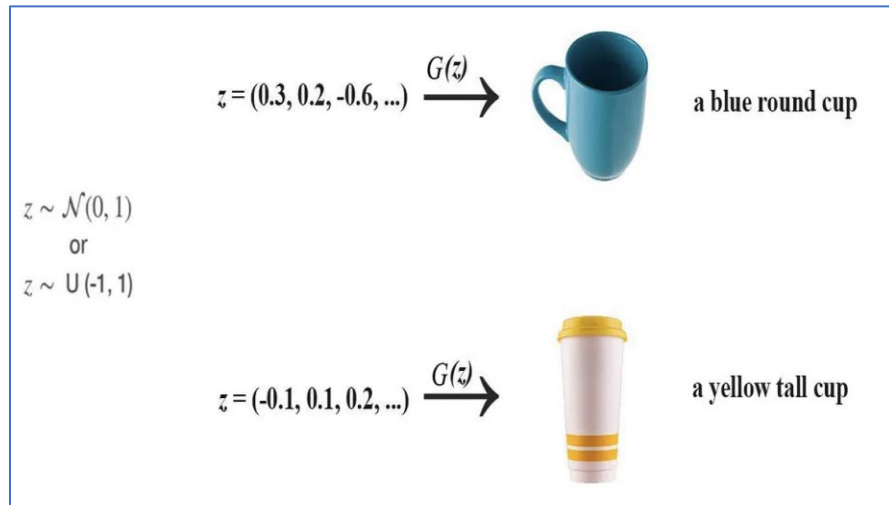
# What's a GAN

- GAN pits two neural networks against each other:
  - a generator network  $G(z)$
  - a discriminator network  $D(x)$ .
- The generator tries to mimic examples from a training dataset, which is sampled from the true data distribution  $q(x)$ . It does so by transforming a random source of noise received as input into a synthetic sample.
- The discriminator receives a sample, but it is not told where the sample comes from. Its job is to predict whether it is a real data sample or a synthetic sample.
- The discriminator is trained to make accurate predictions, and the generator is trained to output samples that fool the discriminator into thinking they came from the data distribution.



# What's G?

- The Generator samples noise  $\mathbf{z}$  using a normal or uniform distribution.
- With  $\mathbf{z}$  as an input, it uses a generator  $\mathbf{G}$  to create an image  $\mathbf{x}$  ( $\mathbf{x}=\mathbf{G}(\mathbf{z})$ )
- It performs multiple transposed de-convolution to upsample  $\mathbf{z}$  and generate the image  $\mathbf{x}$



- The Generator alone will just create random noise.
- The Discriminator “conceptually” provides guidance to the Generator on what images to create.

# Information Theory Background

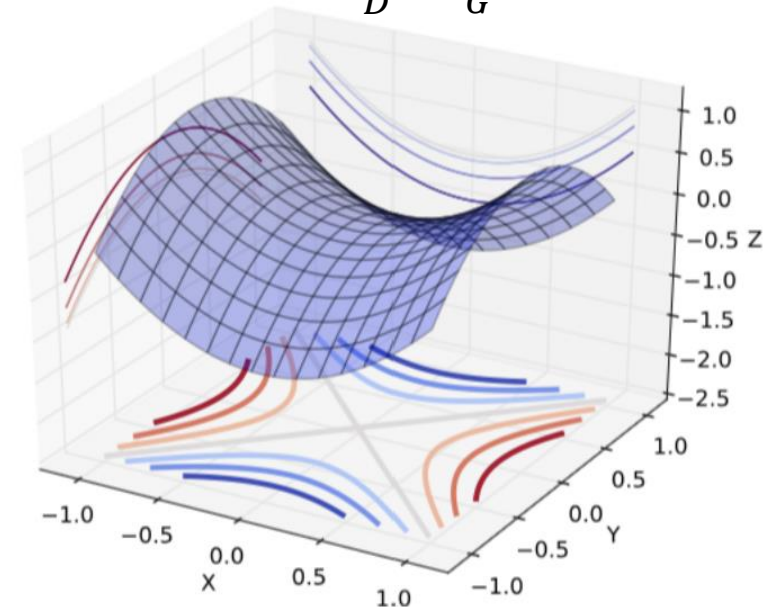
- The information  $I$  of an event  $x$  with probability distribution  $P(x)$  is defined as:  $I(x) = -\log_2 P(x)$
- The entropy  $H$  measures the amount of information and is defined as:  $H(x) = E_{x \sim P}[I(x)]$
- Then we have:  $H(x) = -E_{x \sim P}[\log_2 P(x)] = \sum_x P(x) \log_2 P(x)$
- Example: entropy of a coin launch:  $H(x) = -p(\text{head}) \log_2 p(\text{head}) - p(\text{tail}) \log_2 p(\text{tail}) = -\frac{1}{2} \log_2 \left(\frac{1}{2}\right) = 1$
- The Cross-Entropy  $H$  of an event  $x$  with probability distribution  $P(x)$ , but relative to a different sample  $z$  with probability distribution  $Q(z)$ , is defined as:  $H(x) = \sum_x P(x) \log_2 Q(z)$
- $D$  is a classifier  $X \rightarrow \{0, 1\}$  that tries to distinguish between:
  - a sample from the data distribution ( $D(x) = 1$ , for  $x \sim p_{\text{data}}$ )
  - a sample from the model distribution ( $D(G(z)) = 0$ , for  $z \sim p_{\text{noise}}$ )
- $G$  is a generator  $Z \rightarrow X$  trained to produce samples  $G(z)$  for  $z \sim p_{\text{noise}}$  that are difficult for  $D$  to distinguish from data.

# GAN Value Function (1)

- Consider the value function:  $V(D, G) = E_{x \sim p_{data}} [\log_2(D(x))] + E_{z \sim p_{noise}} [\log_2(1 - D(G(z)))]$
- We want to:
  - for fixed  $G$ , find  $D$  which maximizes  $V(D, G)$
  - for fixed  $D$ , find  $G$  which minimizes  $V(D, G)$
- In other words, we are looking for the saddle point:
- An “optimal” discriminator is :

$$D_G^*(x) = \frac{p_{data}(x)}{p_{data}(x) + p_G(x)}$$

$$(D^*, G^*) = \max_D \min_G V(D, G)$$



# GAN Value Function (2)

- The  $D_{KL}$  Kullback-Leibler divergence is defined as:  $D_{KL}(P||Q) = E_x \log_2 \frac{P(x)}{Q(x)}$
- Using the optimal discriminator  $D_G^*(x)$  we have:

$$D_G^*(x) = \frac{p_{data}(x)}{p_{data}(x) + p_G(x)}$$

$$V(D, G) = D_{KL}\left[p_{data}, \frac{p_{data}(x) + p_G(x)}{2}\right] + D_{KL}\left[p_G, \frac{p_{data}(x) + p_G(x)}{2}\right] = \underbrace{D_{JSD}[p_{data}(x), p_G(x)]}_{2 \times \text{Jensen-Shannon Divergence (JSD)}} - \log_2(4)$$

notice the difference



# The GAN Training Algorithm

- Sample minibatch of  $m$  training points  $x(1), x(2), \dots, x(m)$  from  $D$
- Sample minibatch of  $m$  noise vectors  $z(1), z(2), \dots, z(m)$  from  $p_z$
- Update the discriminator parameters  $\phi$  by stochastic gradient **ascent**:

$$\nabla_{\phi} V(G_{\theta}, D_{\phi}) = \frac{1}{m} \nabla_{\phi} \sum_{i=1}^m [\log D_{\phi}(x^{(i)}) + \log(1 - D_{\phi}(G_{\theta}(z^{(i)})))]$$

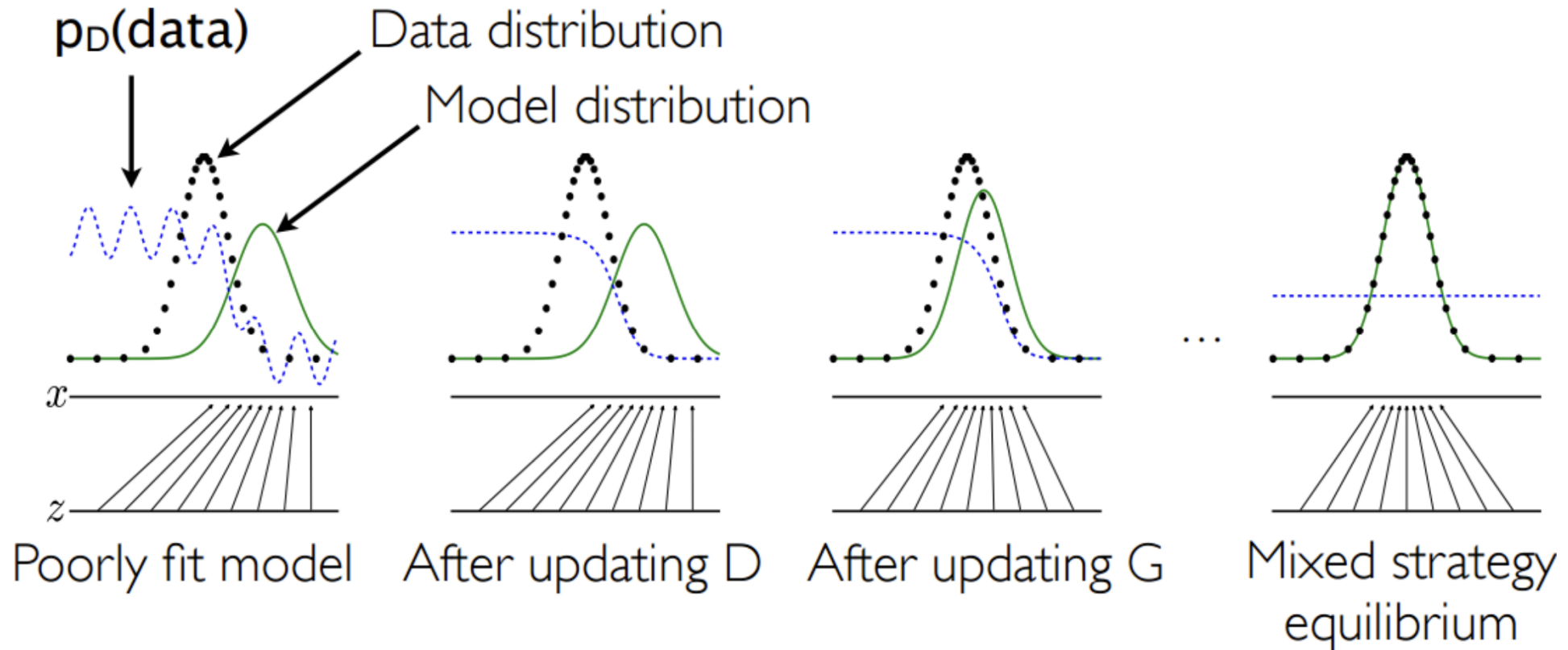
- Update the generator parameters  $\vartheta$  by stochastic gradient **descent**:

$$\nabla_{\theta} V(G_{\theta}, D_{\phi}) = \frac{1}{m} \nabla_{\theta} \sum_{i=1}^m \log(1 - D_{\phi}(G_{\theta}(z^{(i)})))$$

- Repeat for fixed number of epochs

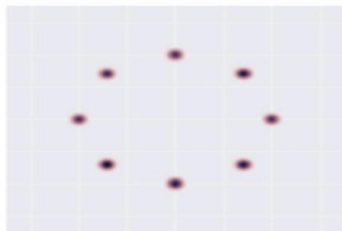
# GAN Gradient Calculation

- Assuming  $D$  and  $G$  are neural networks parameterized by  $\vartheta_D$  and  $\vartheta_G$ , backpropagation can be used to optimize  $D$ 's and  $G$ 's objectives alternatively until convergence.



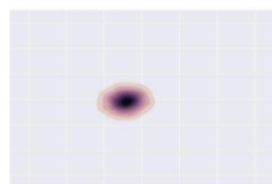
# GANs Collapse

- GANs are notorious for suffering from mode collapse
- Intuitively, this refers to the phenomena where the generator of a GAN collapses to one or few samples (dubbed as “modes”)
- There is not a general solution but just empirical-driven methods (*Q-GAN may have it*)
- Example: the target distribution is a mixture of Gaussians

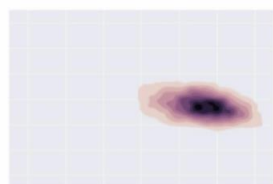


Target

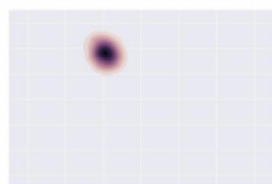
- The generator distribution keeps oscillating between different modes



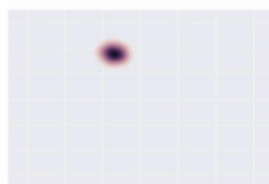
Step 0



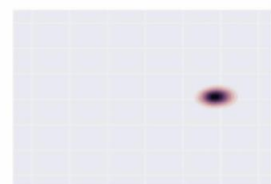
Step 5k



Step 10k



Step 15k



Step 20k

# Demo Classical GAN

**Demo Classical GAN**

# Quantum Computing Short Introduction

- Quantum Logic, Gates and Circuits
- Quantum Algorithms
- Demo Pennylane Quantum GAN
- Demo Qiskit Quantum GAN

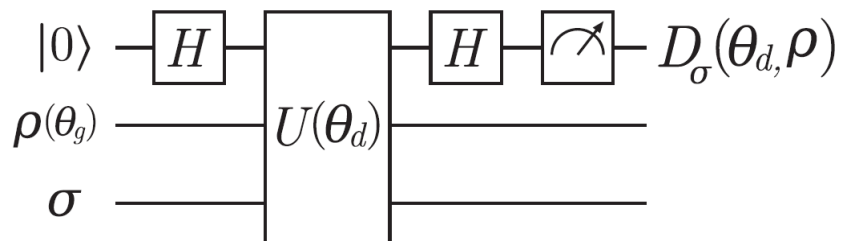
**[Go to Q.C. Intro](#)**

# Quantum GAN

- Quantum GANs are the transposition of the GAN model on quantum computing
- Generalizing their framework to the quantum settings:
- classical data  $\rightarrow$  density matrix  $\sigma = \sum_i p_i |\psi_i\rangle\langle\psi_i|$
- Generator  $\rightarrow$  unitary operator on a quantum circuit  $U$  that outputs the quantum state  $\rho$ :  
 $\rho = U(\theta_g)\rho_0U^\dagger(\theta_g)$  from the initial state  $\rho_0$
- Discriminator  $\rightarrow$  positive operator  $T$  that takes both true data  $\sigma$  and fake data  $\rho$  in input and outputs the probability the data being true:  $D(\theta_d, \rho_{\text{in}}) = \text{Tr}[T\rho_{\text{in}}]$
- The Q-GAN must solve the “minimax game”:  $\min_{\theta_g} \max_T (\text{Tr}[T\sigma] - \text{Tr}[T\rho(\theta_g)])$

# Convergence of Quantum GAN

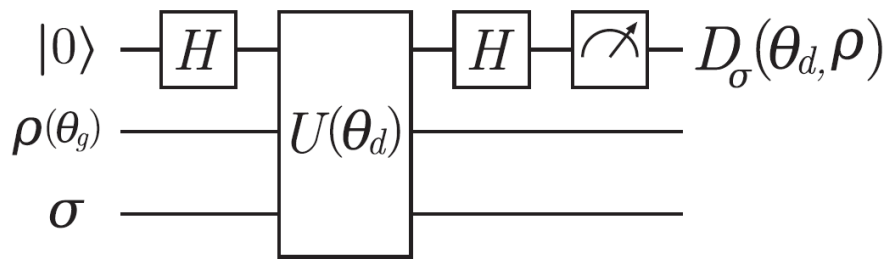
- To evaluate the fidelity of the  $\rho$  density matrix respect to  $\sigma$  the usage of Helstrom measurement operator  $T = P^+(\sigma - \rho)$  is not a good choice since it may lead to oscillation
- Niu et alii proposed to leverage a specific feature of Quantum Computing that is the entanglement: rather than evaluating either fake or true data individually the proposed discriminator perform a measurement on the **joint system** of the true data  $\sigma$  and the generated data  $\rho(\theta_g)$
- The «entangled» fidelity is:  $D_{\sigma}^{\text{fid}}[\rho(\theta_g)] = [\text{Tr} \sqrt{\sigma^{1/2} \rho(\theta_g) \sigma^{1/2}}]^2$



EQ-GAN architecture:

- the generator varies  $\theta_g$  to fool the discriminator;
- the discriminator varies  $\theta_d$  to distinguish the state;
- the global optimum of the EQ-GAN occurs when  $\rho(\theta_g) = \sigma$

# Entangling Q-GAN



Given this architecture the unitary  $U$  operator must be implemented in order to evaluate :

$$\min_{\theta_g} \max_{\theta_d} V(\theta_g, \theta_d) = \min_{\theta_g} \max_{\theta_d} \{1 - D_{\sigma}[\theta_d, \rho(\theta_g)]\}$$

Instead simply swapping  $\rho$  and  $\sigma$ , let's take  $U$  such that:  $U(\theta_d) = \exp[-i\theta_d \text{CSWAP}]$

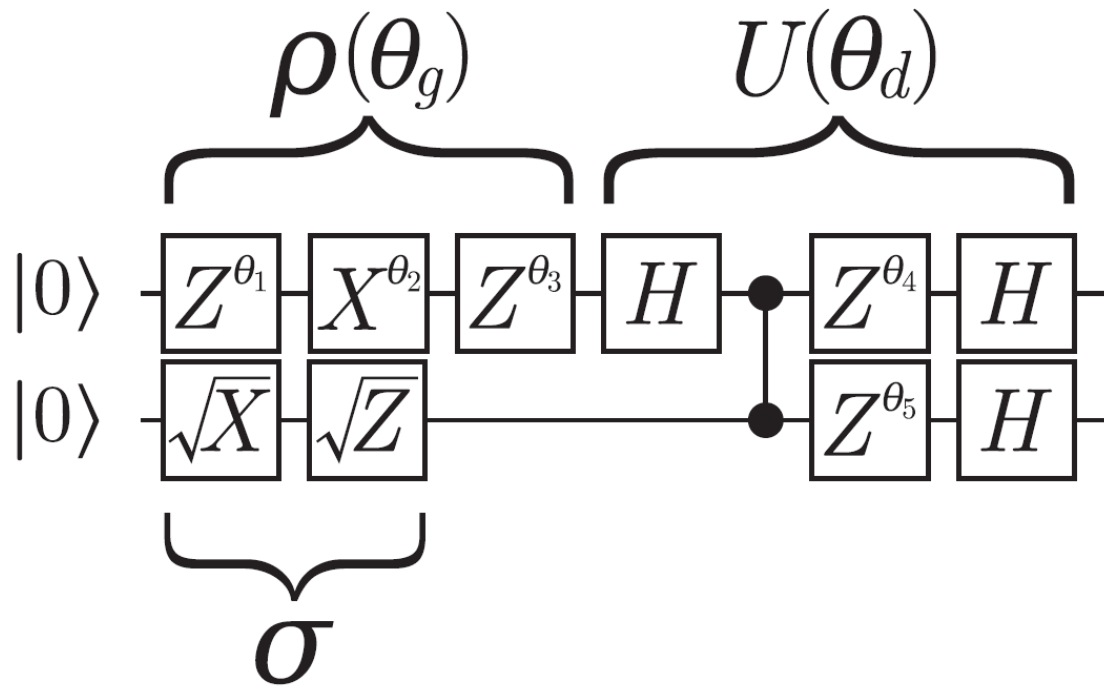
So we have:

$$\begin{aligned} HU(\theta_d)H|0\rangle_a|\psi\rangle|\zeta\rangle &= \frac{i \sin \theta_d}{2} |1\rangle_a [|\zeta\rangle|\psi\rangle - |\psi\rangle|\zeta\rangle] \\ &\quad + \frac{1}{2} |0\rangle_a [(e^{-i\theta_d} + \cos \theta_d)|\psi\rangle|\zeta\rangle \\ &\quad - i \sin \theta_d |\zeta\rangle|\psi\rangle] \end{aligned}$$

$$D_{\sigma}[\theta_d, \rho(\theta_g)] = \frac{1}{2} \{1 + \cos^2 \theta_d + \sin^2 \theta_d D_{\sigma}^{\text{fid}}[\rho(\theta_g)]\}.$$



# Entangling Q-GAN Circuit



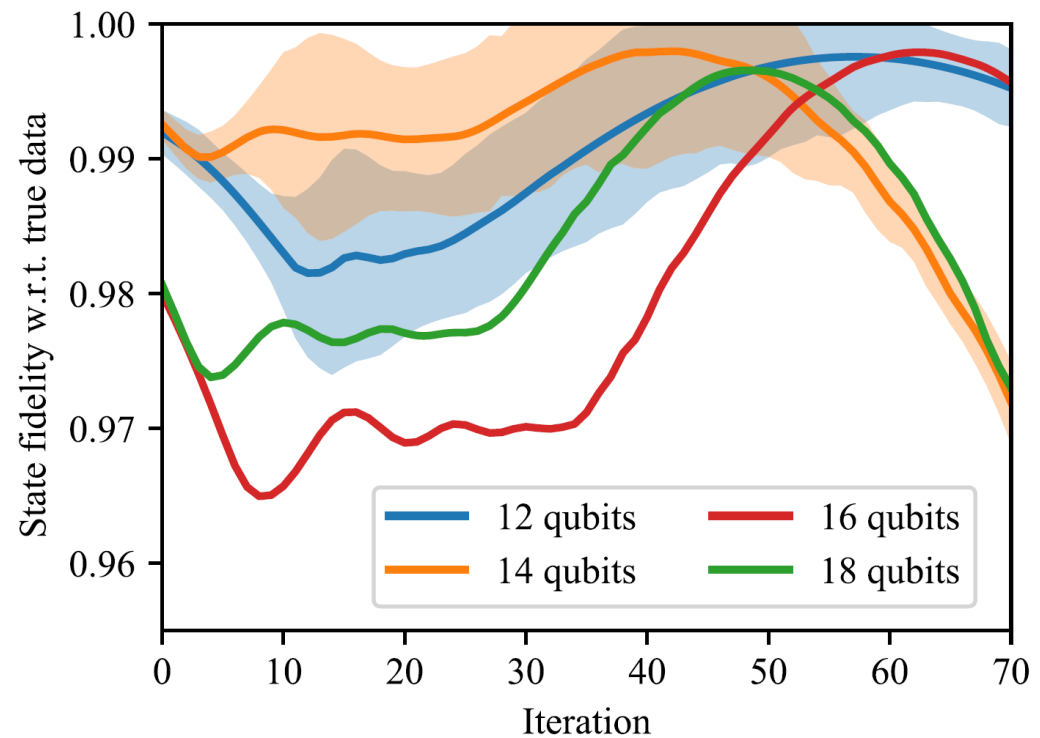
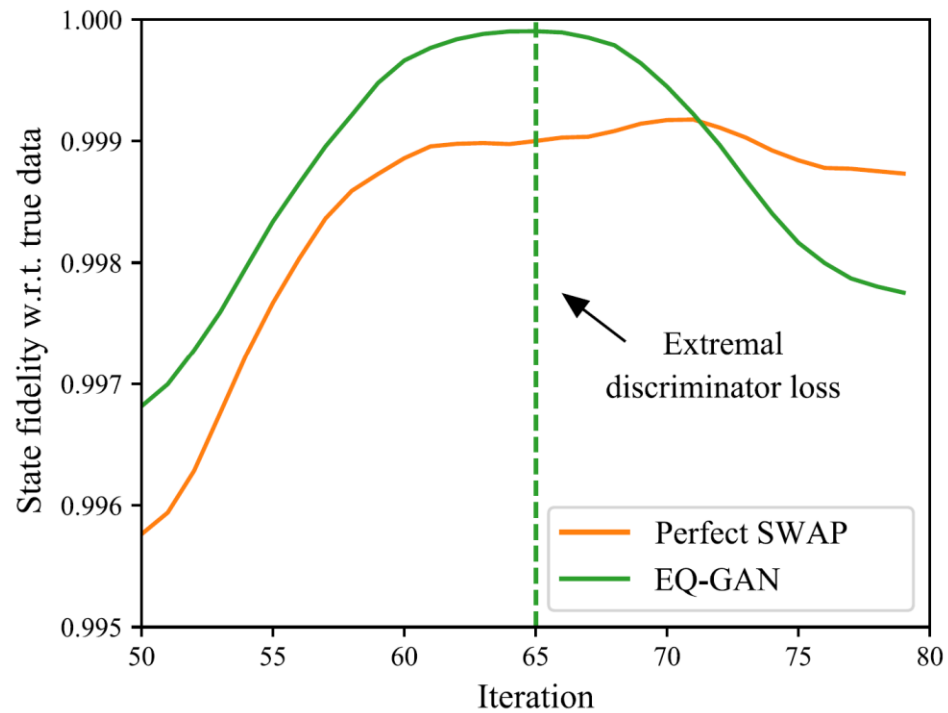
[View the circuit in the IBM composer](#)

$\theta_1, \theta_2, \dots, \theta_5$  are the rotation parameters that are randomly set and then modified by the learning algorithm

# Entangling Q-GAN Results

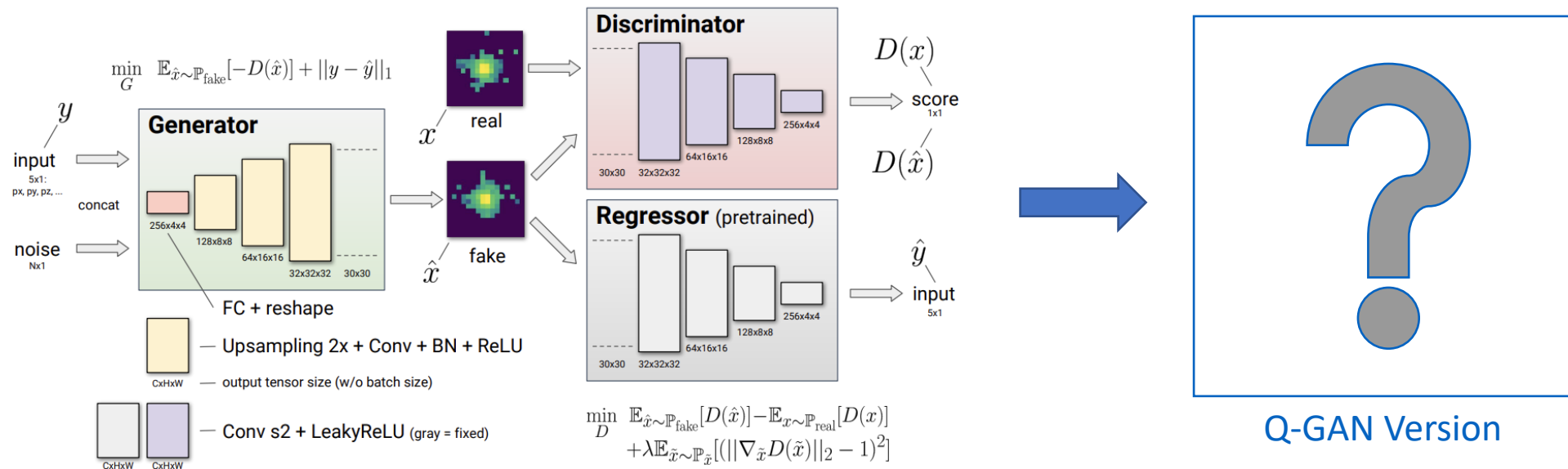
Minimum error in state fidelity  $\rightarrow$  (1-state fidelity)

QML model	Minimum error in state fidelity
Perfect swap	$(2.4 \pm 0.5) \times 10^{-4}$
<b>EQ-GAN</b>	<b><math>(0.6 \pm 0.2) \times 10^{-4}</math></b>



# Q-GAN in Physics

- **Dual-Parameterized Quantum Circuit GAN Model in High Energy Physics**  
Su Yeon Chang, Steven Herbert, Sofia Vallecorsa, Elías F. Combarro, Ross Duncan
- **Quantum Generative Adversarial Networks in a Continuous-Variable Architecture to Simulate High Energy Physics Detectors**  
Su Yeon Chang, Sofia Vallecorsa, Elías F. Combarro, Federico Carminati
- **Generative Adversarial Networks for LHCb Fast Simulation**  
Fedor Ratnikov



# Demo Classical GAN

**Demo Quantum GAN**

# References

- Quantum Computation and Quantum Information  
Michael A.Nielsen – Isaac L.Chuang – Cambridge Press
- An Introduction to Quantum Computing  
Philippe Kaye – Raymond Laflamme – Michele Mosca – Oxford University Press
- Entangling Quantum Generative Adversarial Networks  
Murphy Yuezhen Niu, Alexander Zlokapa, Michael Broughton, Sergio Boixo, Masoud Mohseni, Vadim Smelyanskyi, Hartmut Neven

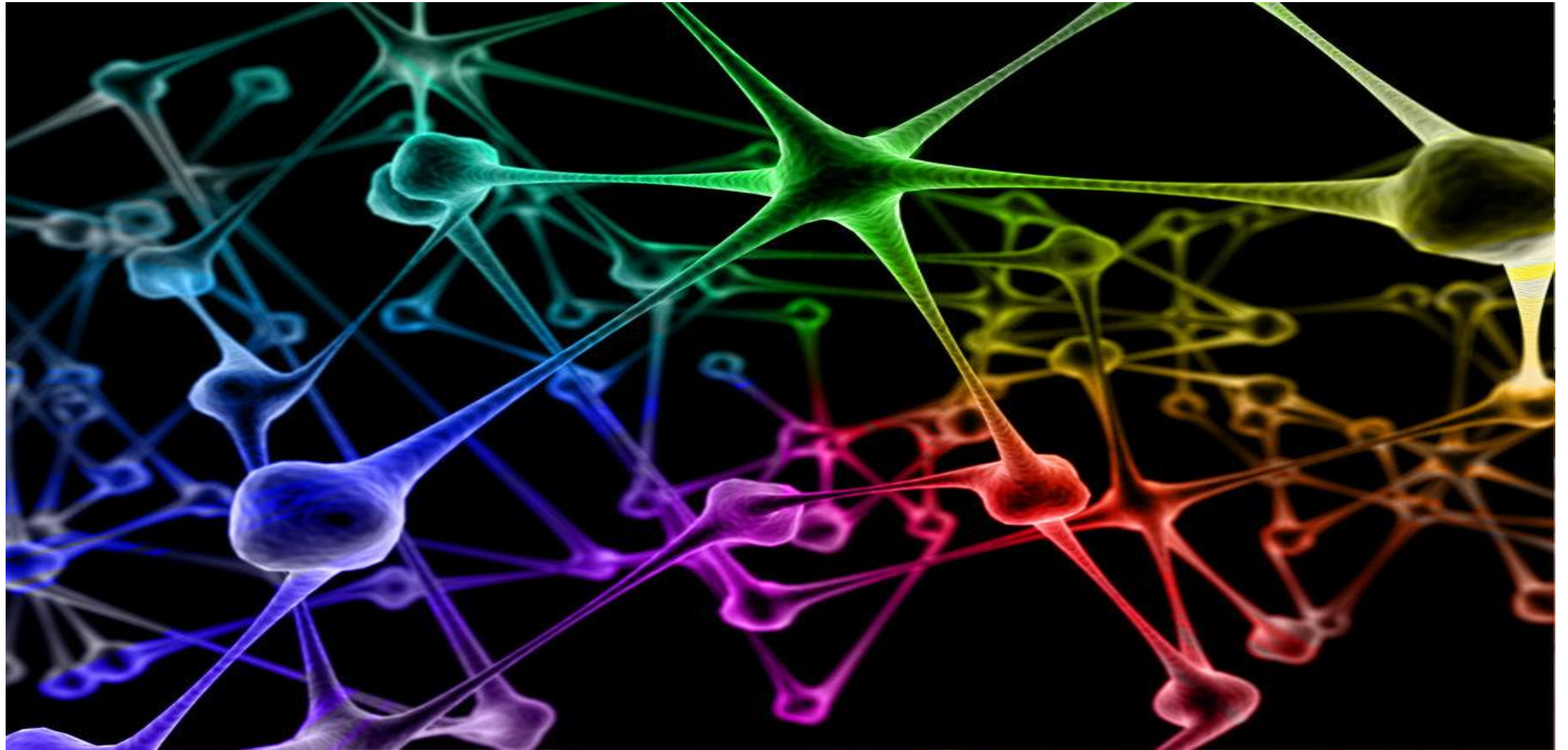


# Questions & Answers

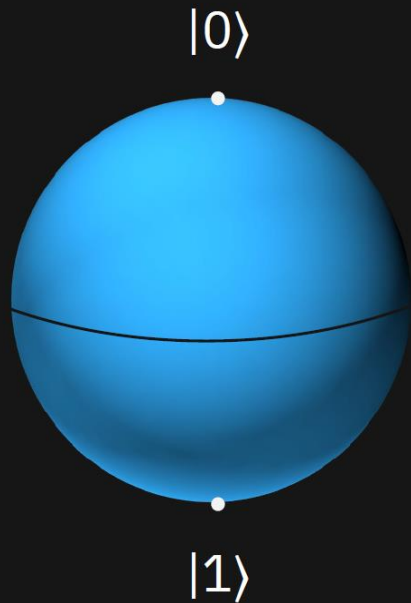


**Q&A**

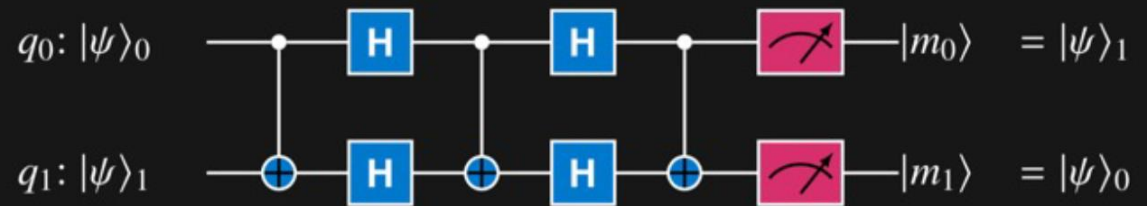
Thanks for your attention!



# Classical vs Quantum Logic



A quantum bit or qubit is a controllable quantum object that is the unit of information



A quantum circuit is a set of quantum gate operations on qubits and is the unit of computation

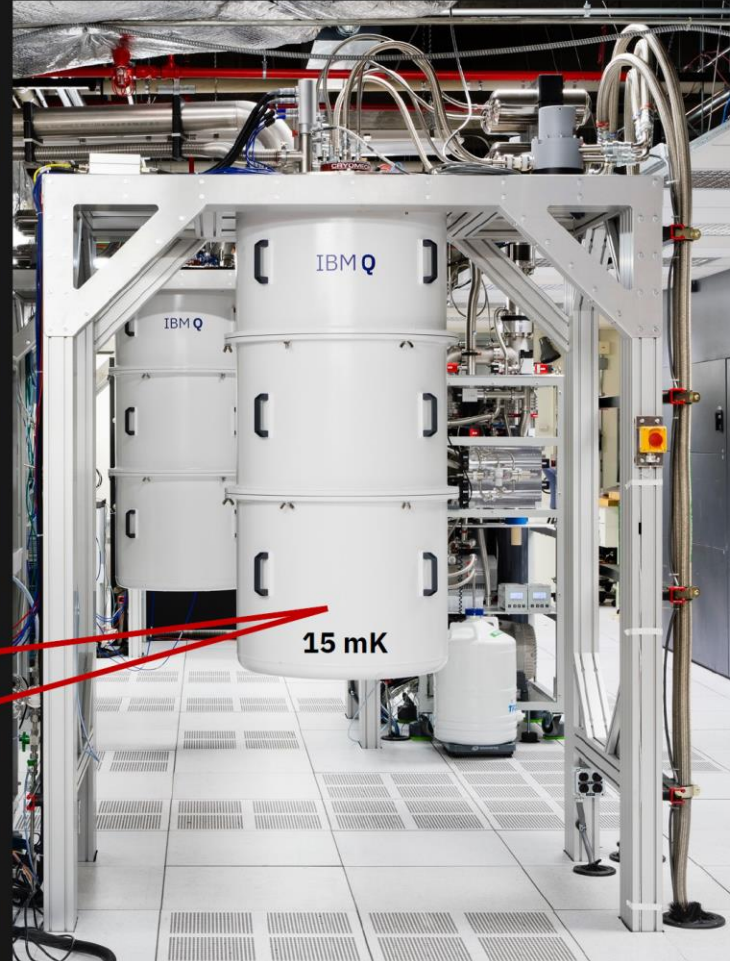
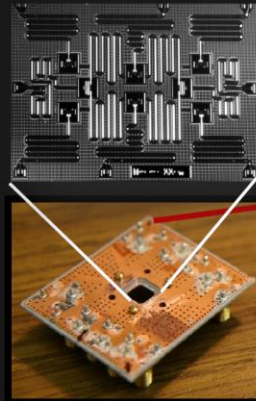


# Quantum Computer Hardware?

## Inside an IBM Quantum system

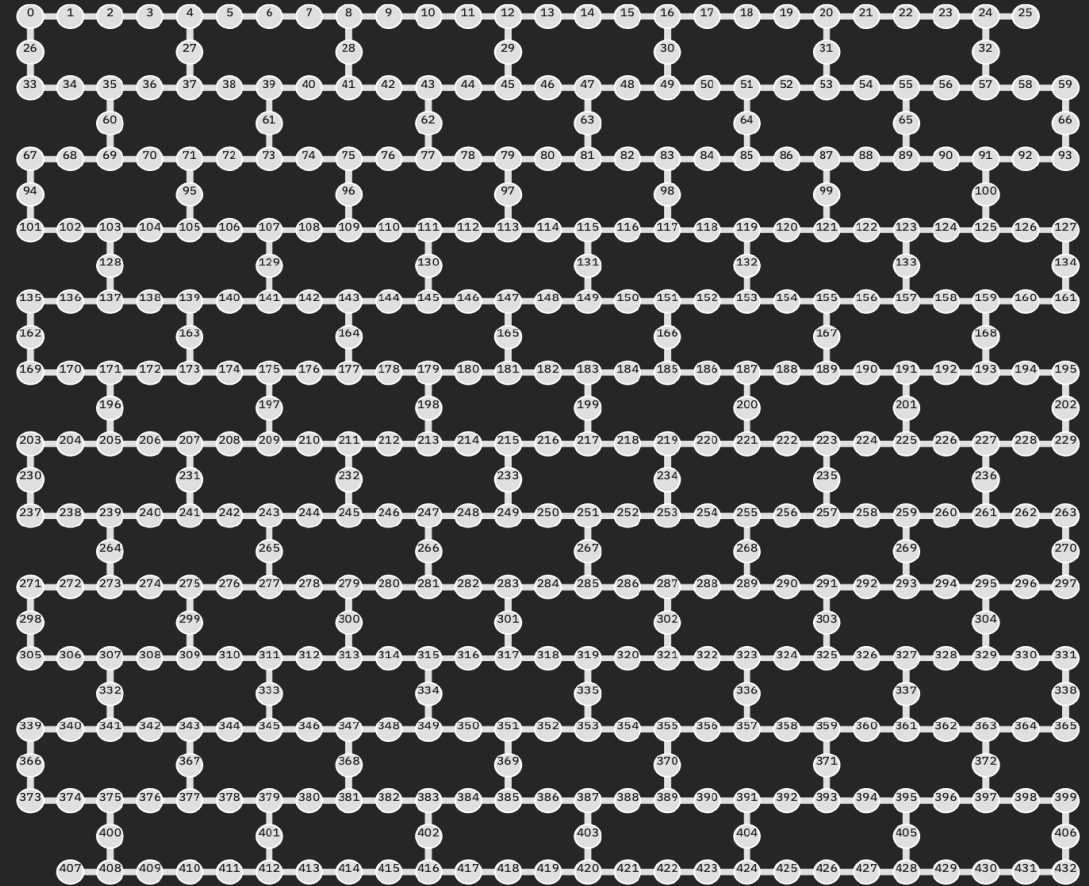
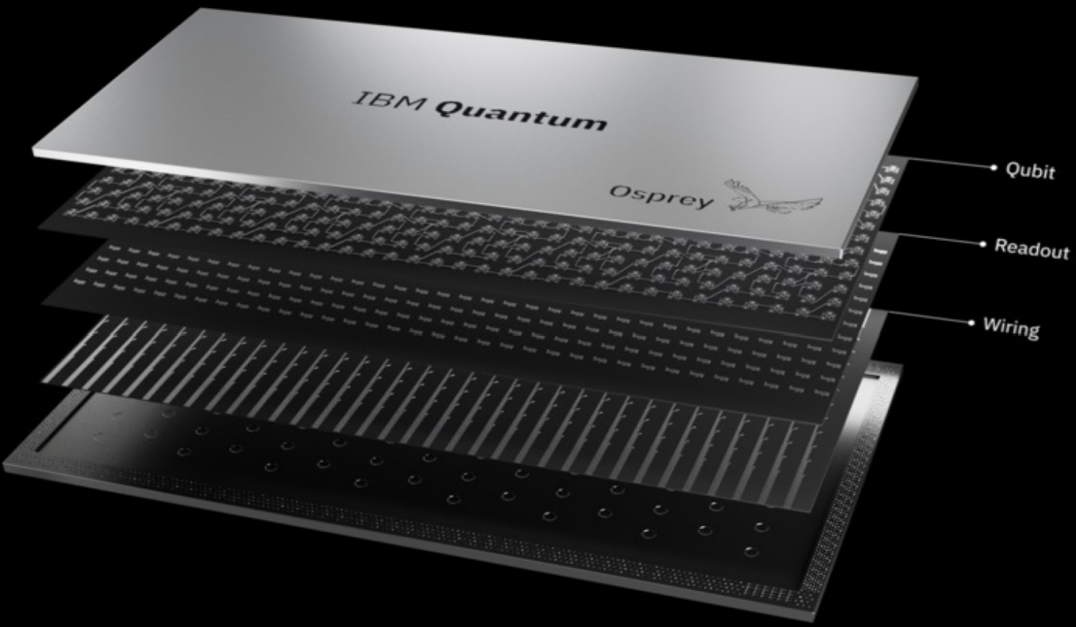
Qubits on chip

Circuit board



# IBM Osprey Chip

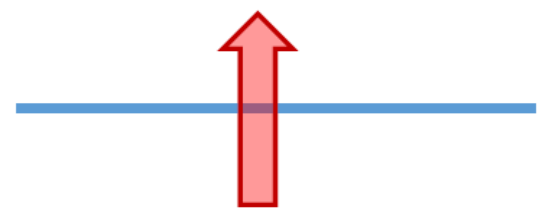
## Osprey – 433 Qubits



Osprey connectivity map

# The basis of Quantum Computing

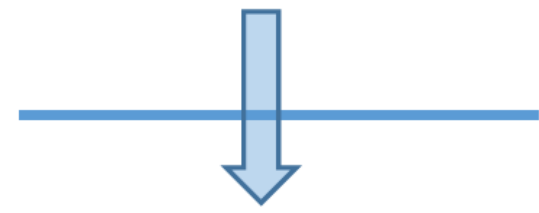
The basis of Quantum Computation is the Qubit:  
a 2-level Quantum Mechanics system



A horizontal blue line representing a qubit state. A red arrow points upwards from the center of the line.

$$|1\rangle = |\uparrow\rangle = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

$$|\psi\rangle = a|0\rangle + b|1\rangle$$



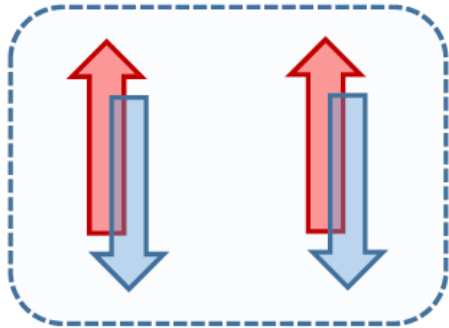
A horizontal blue line representing a qubit state. A blue arrow points downwards from the center of the line.

$$|0\rangle = |\downarrow\rangle = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$$

# What does Quantum Computing rely on?

Qubits can be stacked together:

$$\mathcal{H}_A \otimes \mathcal{H}_B$$



$$|\psi\rangle_A \otimes |\psi\rangle_B \quad \text{basis:} \quad \begin{pmatrix} 1 \\ 0 \end{pmatrix} \otimes \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \begin{pmatrix} 1 \\ 0 \\ 0 \\ 0 \end{pmatrix} ; \quad \begin{pmatrix} 1 \\ 0 \end{pmatrix} \otimes \begin{pmatrix} 0 \\ 1 \end{pmatrix} = \begin{pmatrix} 0 \\ 1 \\ 0 \\ 0 \end{pmatrix} \quad \text{etc.}$$

$$\text{general state: } |\psi\rangle_{AB} = a_{00}|00\rangle + a_{01}|01\rangle + a_{10}|10\rangle + a_{11}|11\rangle$$

$$\text{non-separable state - entangled: } |\psi\rangle_{Bell} = \frac{1}{\sqrt{2}} (|00\rangle + |11\rangle) \neq |\psi\rangle_A \otimes |\psi\rangle_B$$

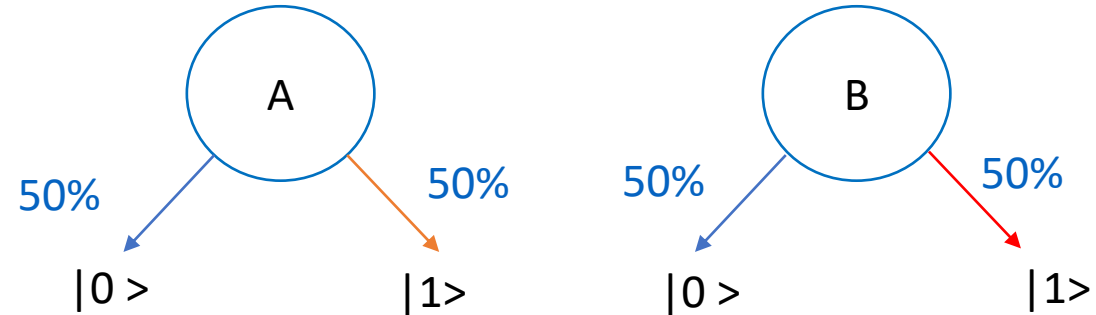
# Entanglement

## Entangled State

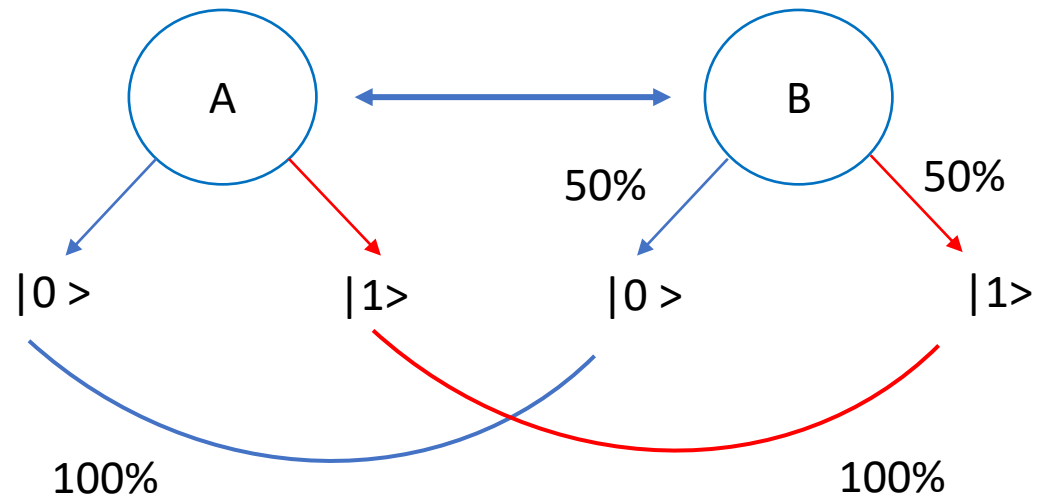
$$|\Phi_+\rangle = \frac{1}{\sqrt{2}}(|00\rangle + |11\rangle)$$

Two entangled states (particles, photons, macroscopic states ...) show a deep correlation even if separated.

## Non-entangled states:



## Entangled states «spooky» action at a distance:



# Important Qubit Gates

$$\text{---} \boxed{X} \text{---} = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \quad \text{“bit flip”}$$

$$\text{---} \boxed{Y} \text{---} = \begin{bmatrix} 0 & -i \\ i & 0 \end{bmatrix} \quad \text{not a straightforward meaning}$$

$$\text{---} \boxed{Z} \text{---} = \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix} \quad \text{“phase flip”}$$

$$|0\rangle \text{---} \boxed{X} \text{---} |1\rangle$$

$$|1\rangle \text{---} \boxed{X} \text{---} |0\rangle$$

“bit flip” is just the classical not gate

Pauli's matrices

$$\begin{aligned} \sigma_0 &= I \\ \sigma_1 &= X \\ \sigma_2 &= Y \\ \sigma_3 &= Z \end{aligned} \quad \sigma_j^2 = I$$

Gate composition

$$\text{---} \boxed{X} \text{---} \boxed{Y} \text{---} = \begin{bmatrix} 0 & -i \\ i & 0 \end{bmatrix} \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} = \begin{bmatrix} -i & 0 \\ 0 & i \end{bmatrix} = \text{---} \boxed{-iZ} \text{---}$$

$$\text{---} \boxed{Y} \text{---} \boxed{X} \text{---} = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} 0 & -i \\ i & 0 \end{bmatrix} = \begin{bmatrix} i & 0 \\ 0 & -i \end{bmatrix} = \text{---} \boxed{iZ} \text{---}$$

# Important Qubit Unitaries

Hadamard gate:  $\text{---} \boxed{H} \text{---} = \begin{bmatrix} \frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} & -\frac{1}{\sqrt{2}} \end{bmatrix}$

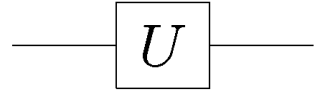
$$H^2 = I$$



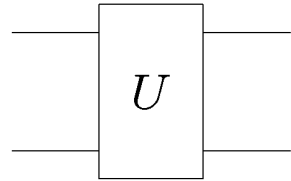
Jacques Hadamard 1865-1963

$$\begin{aligned} \text{---} \boxed{H} \text{---} \boxed{Z} \text{---} \boxed{H} \text{---} &= \begin{bmatrix} \frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} & -\frac{1}{\sqrt{2}} \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix} \begin{bmatrix} \frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} & -\frac{1}{\sqrt{2}} \end{bmatrix} \\ &= \begin{bmatrix} \frac{1}{\sqrt{2}} & -\frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} \end{bmatrix} \begin{bmatrix} \frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} & -\frac{1}{\sqrt{2}} \end{bmatrix} = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} = \text{---} \boxed{X} \text{---} \\ \text{---} \boxed{H} \text{---} \boxed{Z} \text{---} \boxed{H} \text{---} &= \text{---} \boxed{X} \text{---} \end{aligned}$$

# Quantum Circuit Elements



single qubit rotations



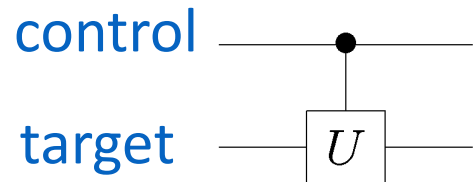
two qubit rotations



controlled-NOT

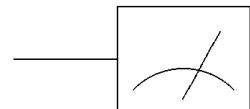
$$= \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}$$

$$\begin{aligned} |00\rangle &\rightarrow |00\rangle \\ |01\rangle &\rightarrow |01\rangle \\ |10\rangle &\rightarrow |11\rangle \\ |11\rangle &\rightarrow |10\rangle \end{aligned}$$



controlled-U

$$= \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & U_{00} & U_{01} \\ 0 & 0 & U_{10} & U_{11} \end{bmatrix}$$



measurement in the  $|0\rangle, |1\rangle$  basis



# Quantum Composer

**Demo Composer**

# Quantum GAN

**Back to Quantum GAN**

# Appendix EQ-GAN

**Appendix EQ-GAN**

# Helstrom Measurement

The Helstrom measurement is the measurement that has the minimum error probability when trying to distinguish between two states.

For example, let's imagine you have two pure states  $|\psi\rangle$  and  $|\phi\rangle$ , and you wish to know which it is that you have. If  $\langle\psi|\phi\rangle = 0$ , then you can specify a measurement with three projectors

$$P_\psi = |\psi\rangle\langle\psi| \quad P_\phi = |\phi\rangle\langle\phi| \quad \bar{P} = \mathbb{I} - P_\psi - P_\phi.$$

(For a two-dimensional Hilbert space,  $\bar{P} = 0$ .)

The question is what measurement should you perform in the case that  $\langle\psi|\phi\rangle \neq 0$ ? Specifically, let's assume that  $\langle\psi|\phi\rangle = \cos(2\theta)$ , and I'll concentrate just on projective measurements (IIRC, this is optimal). In that case, there is always a unitary  $U$  such that

$$U|\psi\rangle = \cos\theta|0\rangle + \sin\theta|1\rangle \quad U|\phi\rangle = \cos\theta|0\rangle - \sin\theta|1\rangle.$$

Now, those states are optimally distinguished by  $|+\rangle\langle+|$  and  $|-\rangle\langle-|$  (you get  $|+\rangle$ , and you assume you had  $U|\psi\rangle$ ). Hence, the optimal measurement is

$$P_\psi = U^\dagger|+\rangle\langle+|U \quad P_\phi = U^\dagger|-\rangle\langle-|U \quad \bar{P} = \mathbb{I} - P_\psi - P_\phi.$$

The success probability is

$$\left(\frac{\cos\theta + \sin\theta}{\sqrt{2}}\right)^2 = \frac{1 + \sin(2\theta)}{2}.$$

More generally, how do you distinguish between two density matrices  $\rho_1$  and  $\rho_2$ ? Start by calculating

$$\delta\rho = \rho_1 - \rho_2,$$

and finding the eigenvalues  $\{\lambda_i\}$  and corresponding eigenvectors  $|\lambda_i\rangle$  of  $\delta\rho$ . You construct 3 measurement operators

$$P_1 = \sum_{i:\lambda_i>0} |\lambda_i\rangle\langle\lambda_i| \quad P_2 = \sum_{i:\lambda_i<0} |\lambda_i\rangle\langle\lambda_i| \quad P_0 = \mathbb{I} - P_1 - P_2.$$

If you get answer  $P_1$ , you assume you had  $\rho_1$ . If you get  $P_2$ , you had  $\rho_2$ , while if you get  $P_0$  you simply guess which you had. You can verify that this reproduces the pure state strategy described above. What's the success probability of this strategy?

$$\frac{1}{2}\text{Tr}((P_1 + P_0/2)\rho_1) + \frac{1}{2}\text{Tr}((P_2 + P_0/2)\rho_2)$$

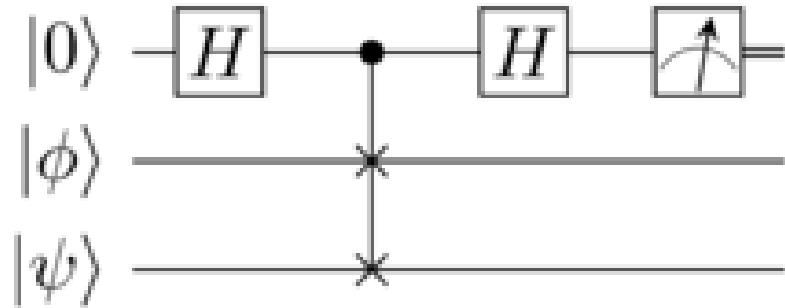
We can expand this as

$$\frac{1}{4}\text{Tr}((P_1 + P_2 + P_0)(\rho_1 + \rho_2)) + \frac{1}{4}\text{Tr}((P_1 - P_2)(\rho_1 - \rho_2))$$

Since  $P_1 + P_2 + P_0 = \mathbb{I}$  and  $\text{Tr}(\rho_1) = \text{Tr}(\rho_2) = 1$ , this is just

$$\frac{1}{2} + \frac{1}{4}\text{Tr}((P_1 - P_2)(\rho_1 - \rho_2)) = \frac{1}{2} + \frac{1}{4}\text{Tr}|\rho_1 - \rho_2|.$$

# Swap Test



- The swap test is a procedure in quantum computation that is used to check how much two quantum states differ, appearing first in the work of Barenco et al.
- It takes two input states and outputs 1 with probability  $\frac{1}{2} - \frac{1}{2} |\langle \psi | \phi \rangle|^2$

# Swap Test

Consider two states:  $|\phi\rangle$  and  $|\psi\rangle$ . The state of the system at the beginning of the protocol is  $|0, \phi, \psi\rangle$ . After the **Hadamard gate**, the state of the system is  $\frac{1}{\sqrt{2}}(|0, \phi, \psi\rangle + |1, \phi, \psi\rangle)$ . The **controlled SWAP gate** transforms the state into  $\frac{1}{\sqrt{2}}(|0, \phi, \psi\rangle + |1, \psi, \phi\rangle)$ . The second

Hadamard gate results in

$$\frac{1}{2}(|0, \phi, \psi\rangle + |1, \phi, \psi\rangle + |0, \psi, \phi\rangle - |1, \psi, \phi\rangle) = \frac{1}{2}|0\rangle(|\phi, \psi\rangle + |\psi, \phi\rangle) + \frac{1}{2}|1\rangle(|\phi, \psi\rangle - |\psi, \phi\rangle)$$

The **measurement gate** on the first qubit ensures that it's 0 with a probability of

$$P(\text{First qubit} = 0) = \frac{1}{2} \left( \langle\phi|\langle\psi| + \langle\psi|\langle\phi| \right) \frac{1}{2} \left( |\phi\rangle|\psi\rangle + |\psi\rangle|\phi\rangle \right) = \frac{1}{2} + \frac{1}{2} |\langle\psi|\phi\rangle|^2$$





when measured. If  $\psi$  and  $\phi$  are **orthogonal** ( $|\langle\psi|\phi\rangle|^2 = 0$ ), then the probability that 0 is measured is  $\frac{1}{2}$ . If the states are equal ( $|\langle\psi|\phi\rangle|^2 = 1$ ), then the probability that 0 is measured is 1.

In general, for  $P$  trials of the swap test using  $P$  copies of  $|\phi\rangle$  and  $P$  copies of  $|\psi\rangle$ , the fraction of measurements that are zero is  $1 - \frac{1}{P} \sum_{i=1}^P M_i$ , so by taking  $P \rightarrow \infty$ , one can get arbitrary precision of this value.

# Appendix Quantum Computing

**Appendix Q.C.**

# Main Algorithms as «Lego Bricks»

- Teleportation
  - Superdense Coding
- 
- Quantum Cryptography
- 
- Deutsch
  - Deutsch-Jozsa
- 
- «Oracle» Problems  
(guess the function)
- 
- Grover
- 
- Search Speed-up
- 
- Quantum Phase Estimation (QPE)
- 
- Linear Algebra / Shor



# Quantum Teleportation

## Alice wants to send her qubit to Bob

She does not know the wave function of her qubit



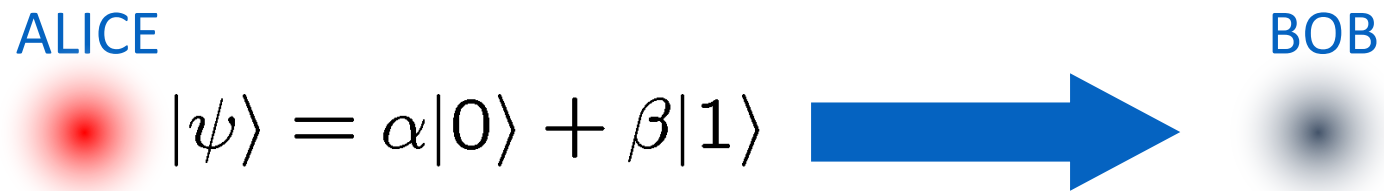
- Alice and Bob share an entangled pair of qubits
- Alice makes a measure on her state and send him classical bits of information about the outcome
- Then Bob would have information about  $|\psi\rangle$  as well

**This would be a procedure for extracting information from without effecting the state  $|\psi\rangle = \alpha|0\rangle + \beta|1\rangle$**

# Quantum Teleportation

Alice shares with Bob an entangled state  $\frac{1}{\sqrt{2}}(|00\rangle + |11\rangle)$

Alice does not know the wave function

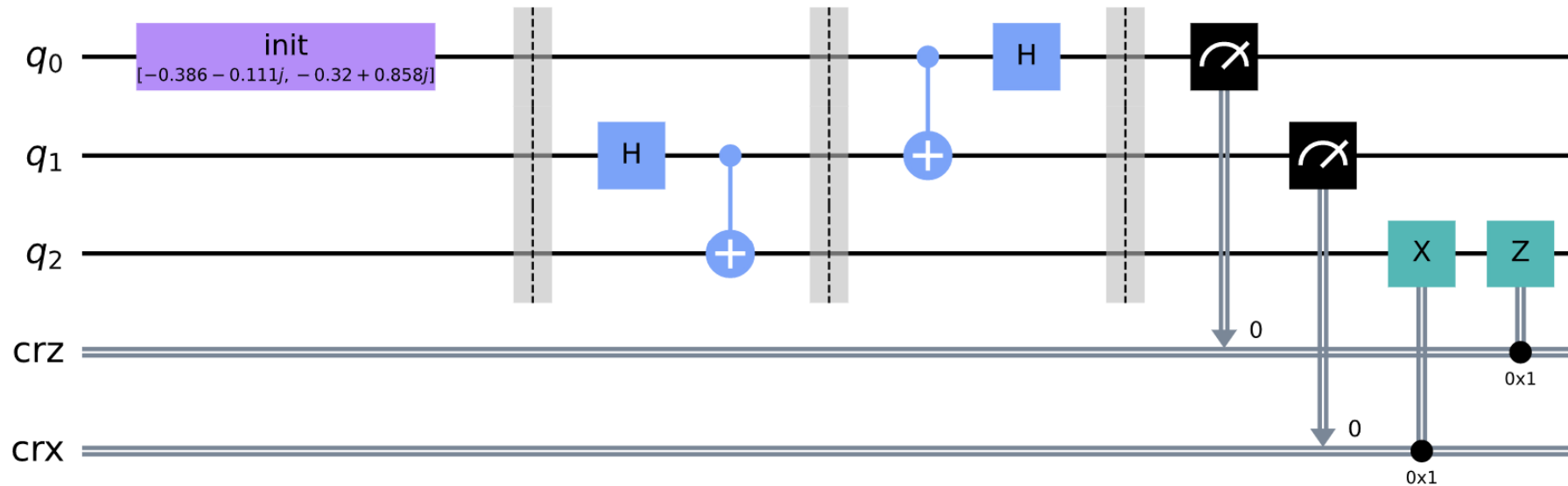


So we have three qubits whose wave function is:

$$\underbrace{|\psi\rangle}_{\text{qubit 1}} \otimes \underbrace{\frac{1}{\sqrt{2}}(|00\rangle + |11\rangle)}_{\text{qubit 2 and qubit 3}}$$

# Quantum Teleportation

Quantum Circuit used to simulate Teleportation:



Alice makes a measurement on  $q_0$  and  $q_1$  and send the bits to Bob

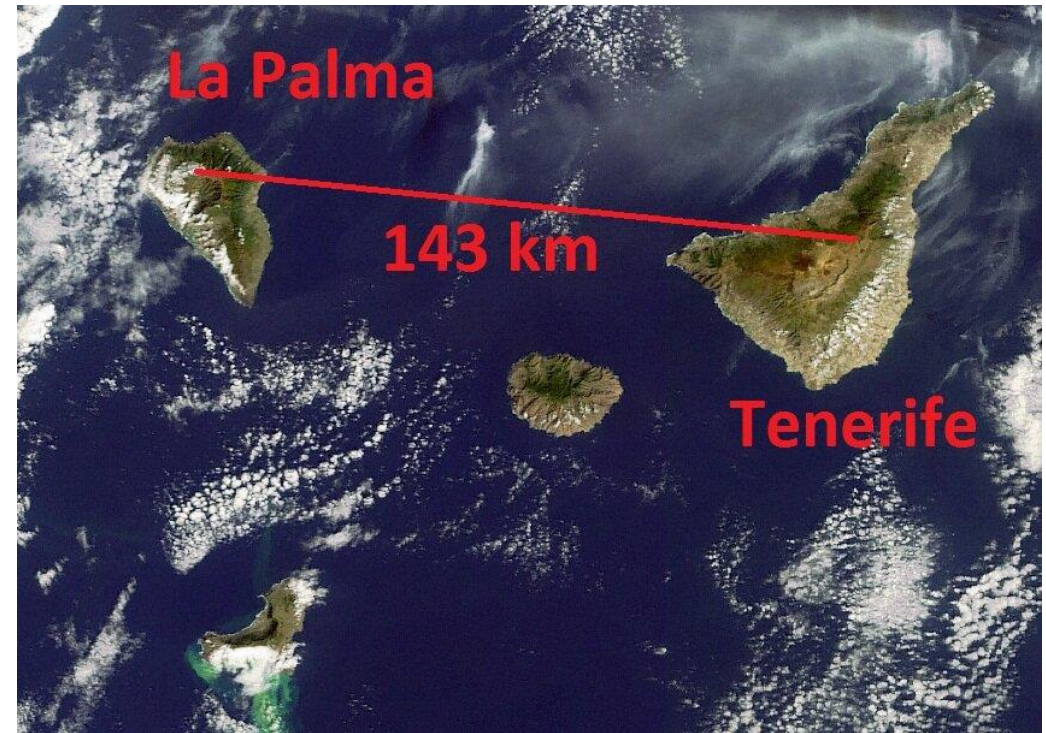
Bob reads them and if

}	00 → does nothing	}	then $q_2 \longleftrightarrow q_0$
	01 → applies X gate		
	10 → applies Z gate		
	11 → applies ZX gate		

# Quantum Teleportation

## ESA observatory breaks world quantum teleportation record

An international research team using ESA's Optical Ground Station in the Canary Islands has set a new distance world record in 'quantum teleportation' by reproducing the characteristics of a light particle across 143 km of open air.



# Quantum Teleportation

**Demo Teleportation**

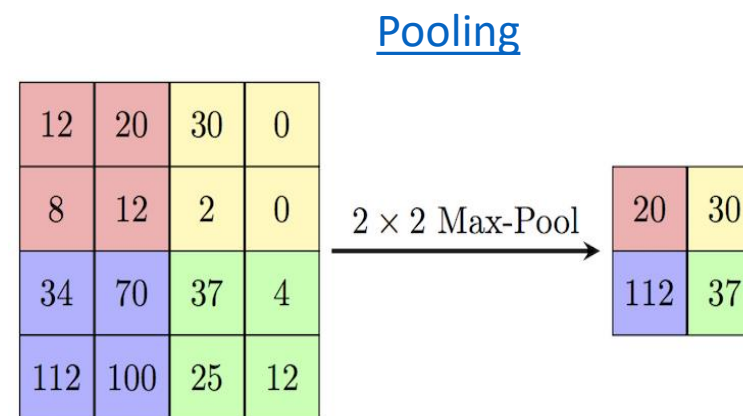
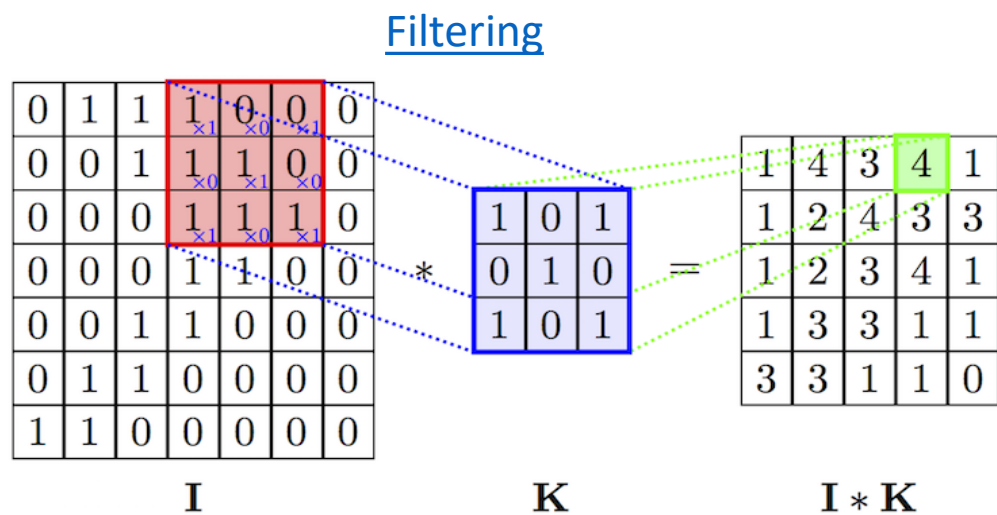
# The Frontier: CNN vs Quantum Machine Learning

- Convolutional neural networks (ConvNets or CNNs) are a category of neural networks<sup>1</sup> that have proven to be very effective in tasks such as image recognition and classification
- Convolution works by extracting a feature map from a source image:



<sup>1</sup><https://ujjwalkarn.me/2016/08/11/intuitive-explanation-convnets/>

# CNN Schema



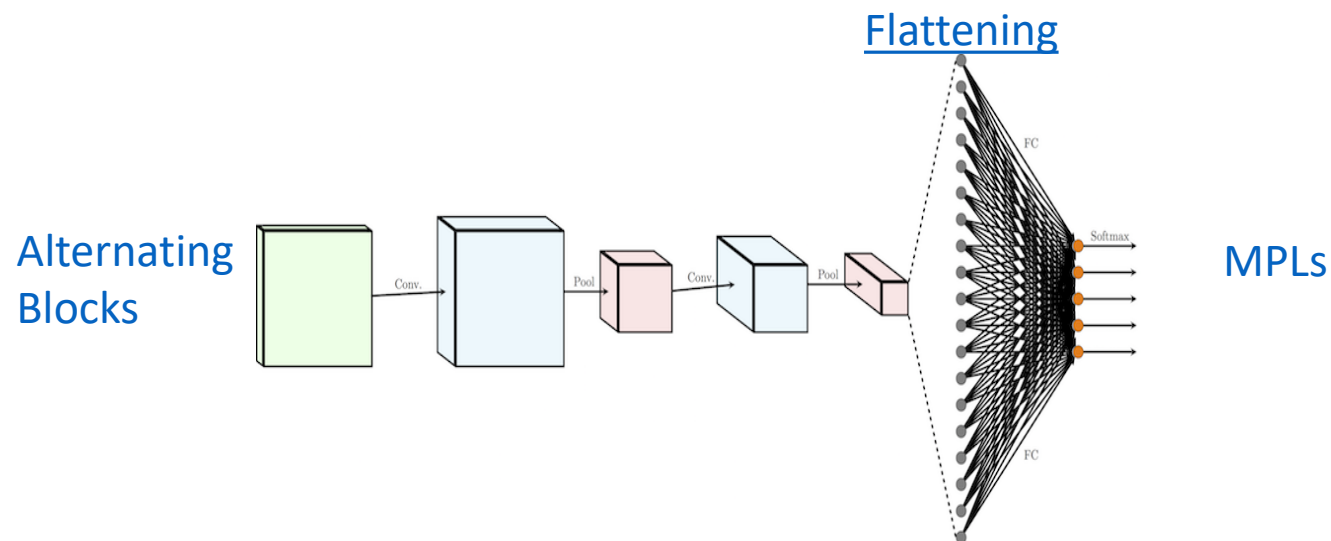
- Convolutional Blocks:

- Filtering
- Pooling

- Alternating Convolutional Blocks

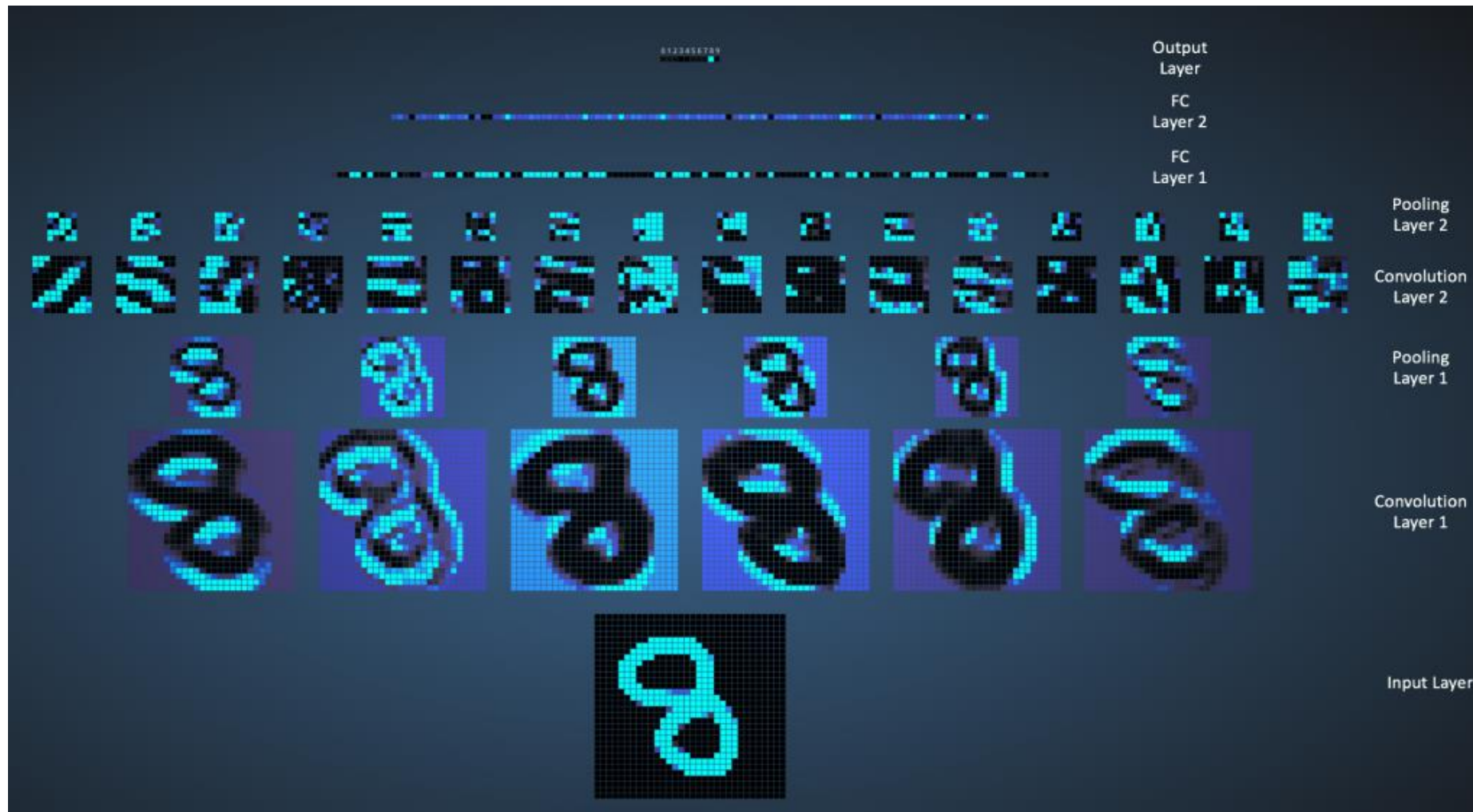
- Flattening

- Multi Perceptron Layers (MPLs)



# CNN ON MNIST

How a convolutional network classifies handwritten numbers:





# CNN AT WORK

**Demo CNN**

# The Frontier: Quantum Machine Learning



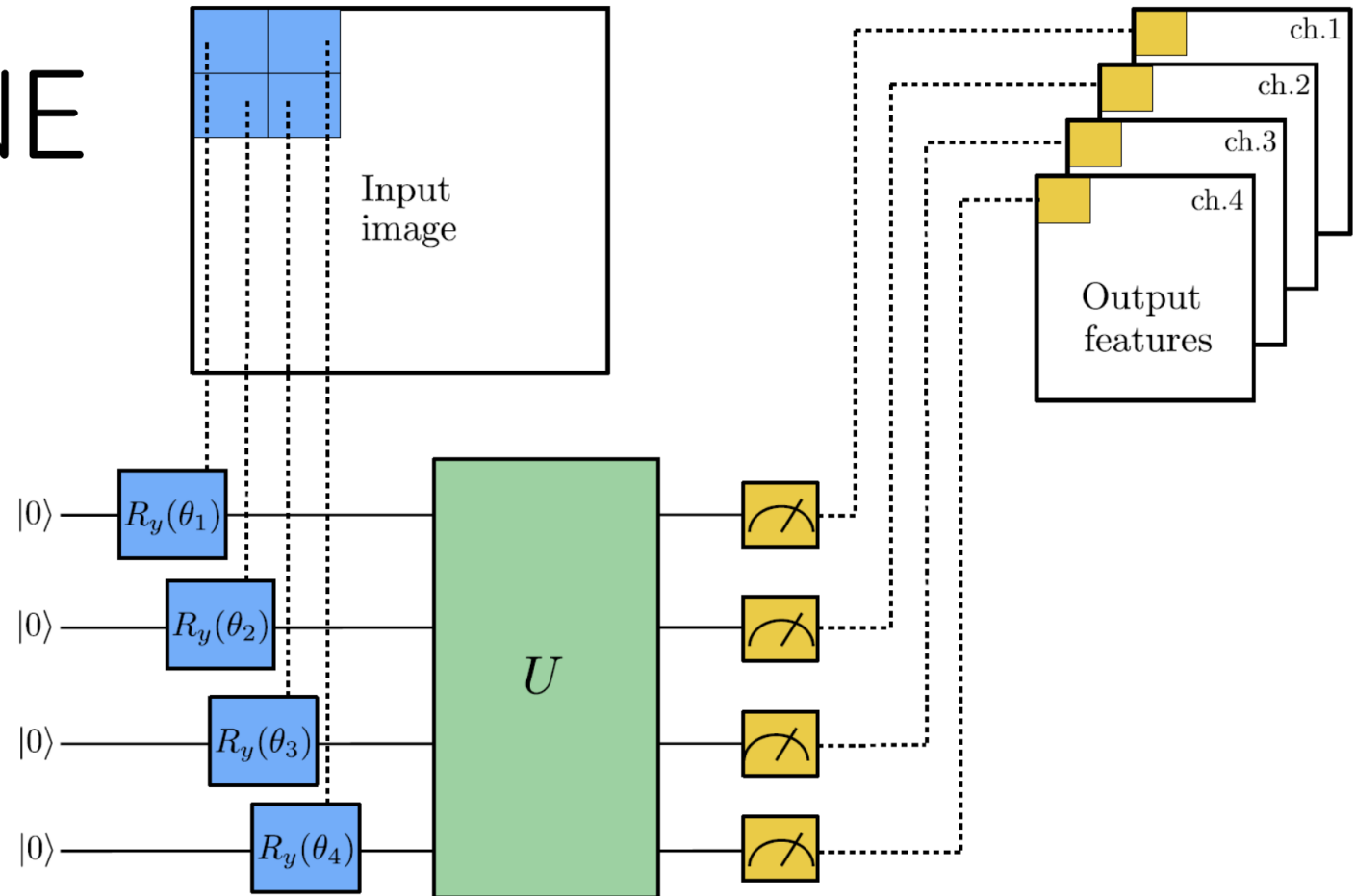
PENNYLANE

## Quantum Machine Learning

Author: Andrea Mari – Posted: 24 March 2020. Last updated: 15 January 2021

Originally introduced by Henderson et alii in 2019

### Pennylane Quantvolution



# Pennylane Quanvolution

**Demo Quanvolution**

Last but not Least...

Let's ask ChatGPT to generate  
qiskit code to demonstrate  
teleportation protocol...

ChatGPT → Qiskit

**Demo ChatGPT**

# More Algorithms

**More Algorithms**

# «Oracle» Problems: Deutsch Algorithm

Given one bit function:  $f : \{0, 1\} \rightarrow \{0, 1\}$

Four such function, it may be:

$$\begin{array}{l} f(0) = f(1) = 0 \\ f(0) = f(1) = 1 \end{array} \left. \vphantom{\begin{array}{l} f(0) = f(1) = 0 \\ f(0) = f(1) = 1 \end{array}} \right\} \text{“constant”}$$
$$\begin{array}{l} f(0) = 0, f(1) = 1 \\ f(0) = 1, f(1) = 0 \end{array} \left. \vphantom{\begin{array}{l} f(0) = 0, f(1) = 1 \\ f(0) = 1, f(1) = 0 \end{array}} \right\} \text{“balanced”}$$



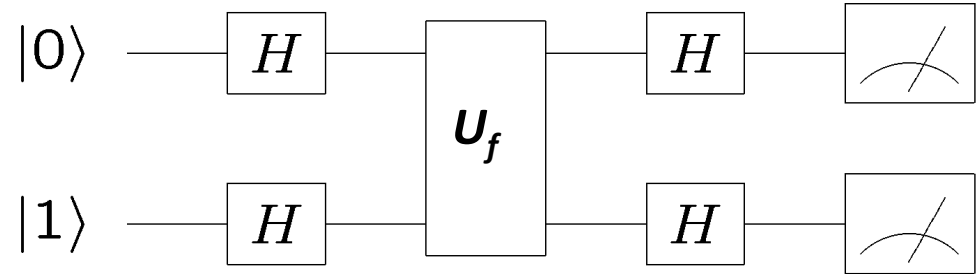
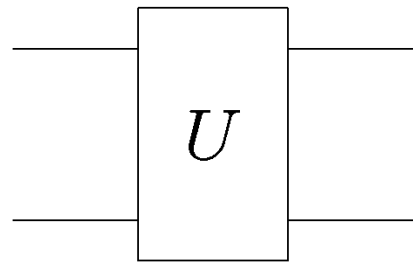
David Deutsch 1985

## Deutsch's Problem

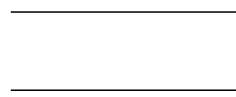
**instance:** function  $f$  is unknown

**problem:** determine whether function is constant or balanced

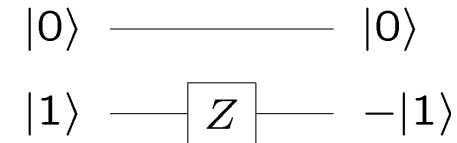
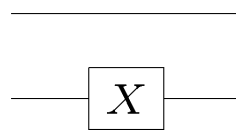
# Quantum Deutsch



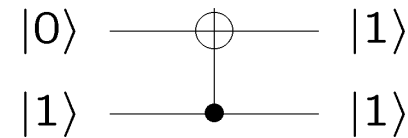
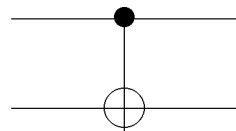
f constant f=0



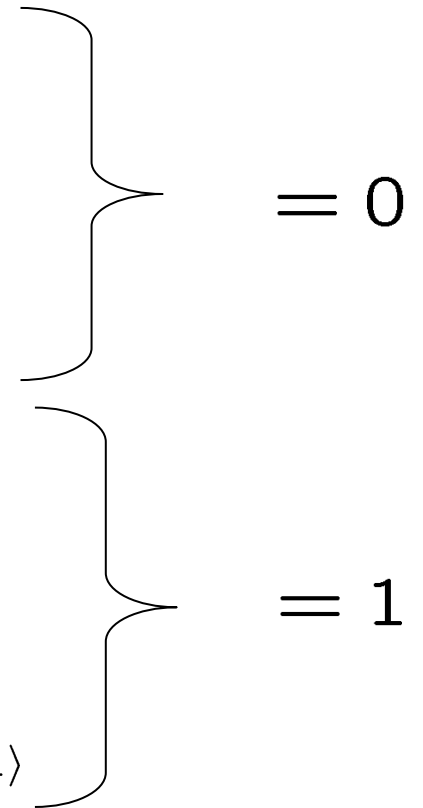
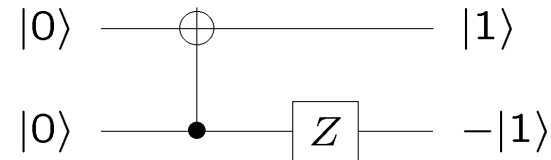
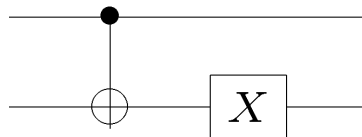
f constant f=1



f balanced f={0,1}



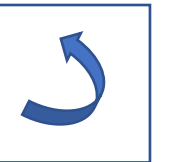
f balanced f={1,0}





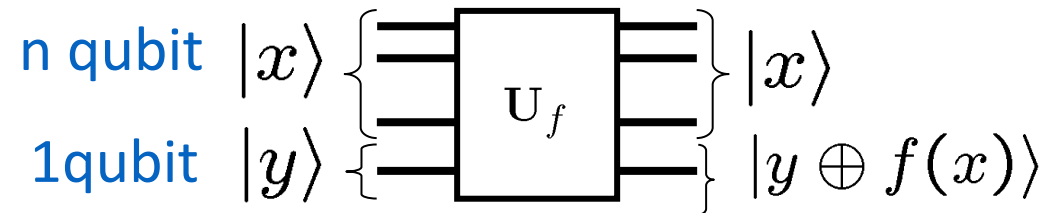
# Deutsch Algorithm

Demo Deutsch



# Searching speed-up: Grover Algorithm

Suppose we have a black box:



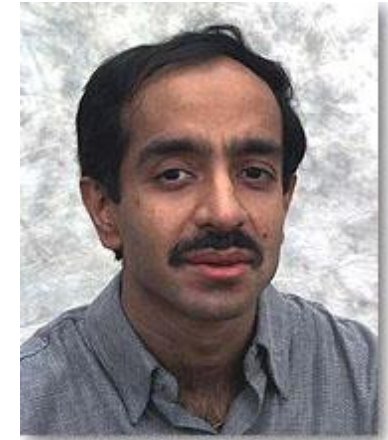
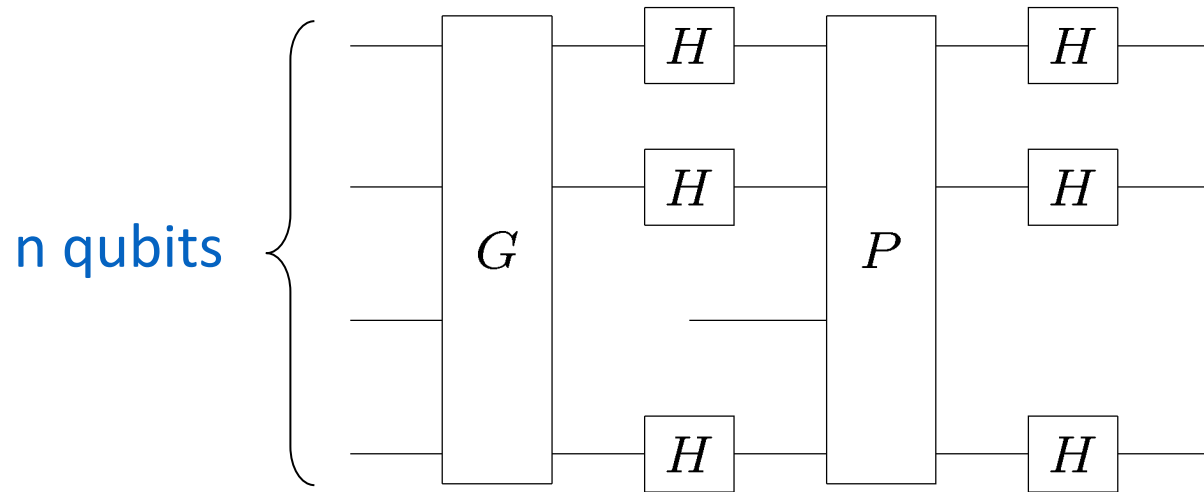
with the property:  $f(x_0) = 1$   
 $\forall x \neq x_0, f(x) = 0$

**Problem:** find  $x_0$  with as few queries as possible

**Solution:** identify the item in “only”  $O(\sqrt{2^n})$  queries!



# Grover Iterate Block



Lov Grover 1996

$$P = \sum_{x=0}^{2^n} -(-1)^{\delta_{x,0}} |x\rangle\langle x| = 2|0\rangle\langle 0| - I$$

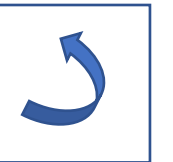
$$H^{\otimes n} P H^{\otimes n} = H^{\otimes n} |0\rangle\langle 0| H^{\otimes n} - I = 2|\psi\rangle\langle\psi| - I$$

$$|\psi\rangle = \frac{1}{\sqrt{2^n}} \sum_{x=0}^{2^n-1} |x\rangle$$

Grover's iterate  $R = (2|\psi\rangle\langle\psi| - I)G$

# Grover Algorithm

**Demo Grover**



# Quantum Phase Estimation Algorithm

We want to estimate a the phase  $\omega$  given the quantum state  $\sum_{y=0}^{2^n-1} e^{2\pi i \omega y} |y\rangle$

Note that in binary we can express  $\omega = 0.x_1x_2x_3 \dots$  as:

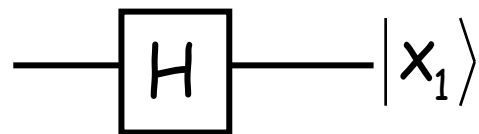
$$2\omega = x_1.x_2x_3 \dots$$

$$2^{n-1}\omega = x_1x_2x_3 \dots x_{n-1}.x_nx_{n+1} \dots$$

If  $\omega=0.x_1$  then we can do the following:

$$\frac{|0\rangle + e^{2\pi i(0.x_1)}|1\rangle}{\sqrt{2}}$$

$$= \frac{|0\rangle + (-1)^{x_1}|1\rangle}{\sqrt{2}}$$



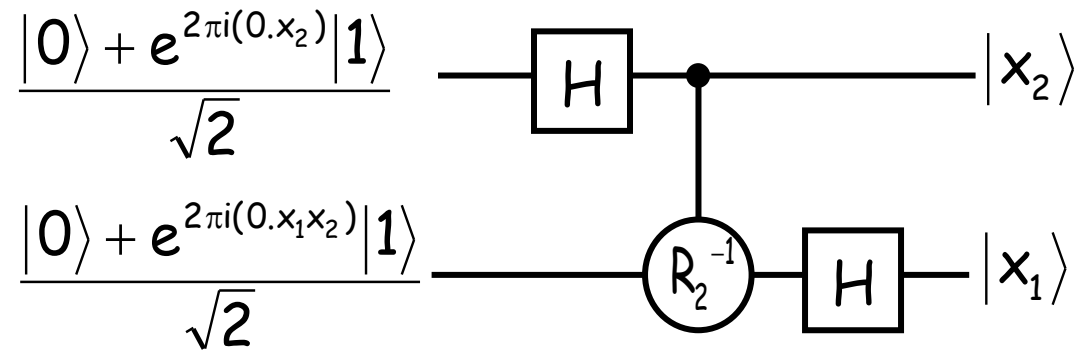
# Quantum Phase Estimation Algorithm

It can be shown that:

$$\sum_{y=0}^{2^n-1} e^{2\pi i \omega y} |y\rangle = \left( |0\rangle + e^{2\pi i (2^{n-1} \omega)} |1\rangle \right) \otimes \left( |0\rangle + e^{2\pi i (2^{n-2} \omega)} |1\rangle \right) \otimes \dots \otimes \left( |0\rangle + e^{2\pi i (\omega)} |1\rangle \right)$$

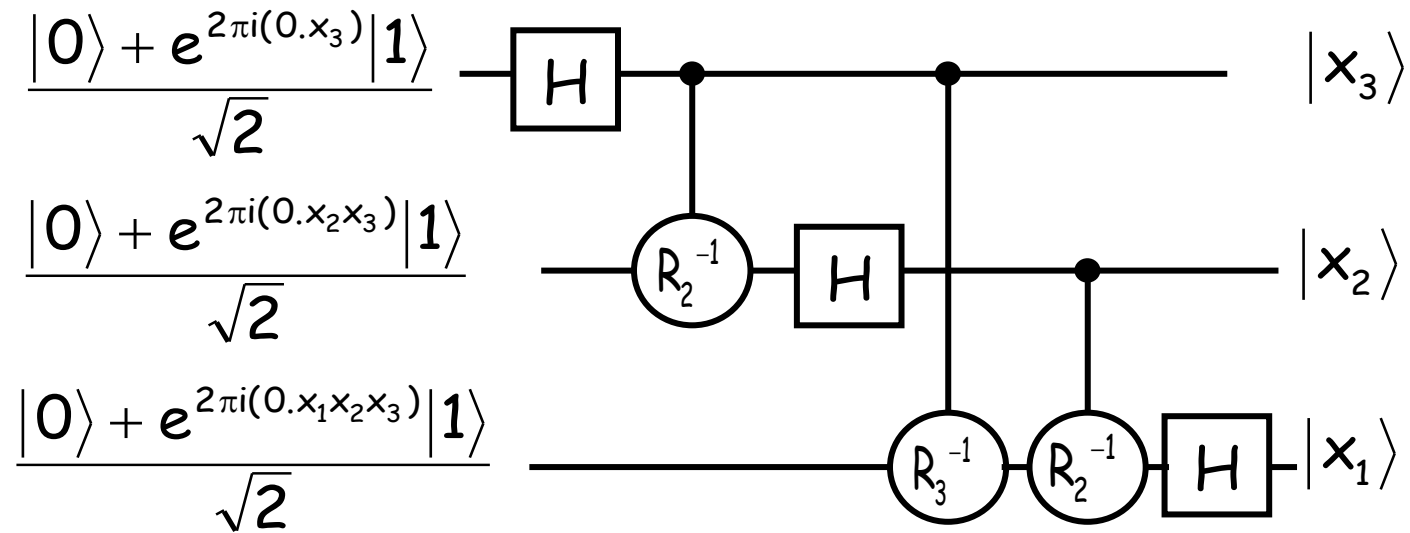
$$= \left( |0\rangle + e^{2\pi i (0.x_n x_{n+1} \dots)} |1\rangle \right) \otimes \left( |0\rangle + e^{2\pi i (0.x_{n-1} x_n x_{n+1} \dots)} |1\rangle \right) \otimes \dots \otimes \left( |0\rangle + e^{2\pi i (0.x_1 x_2 \dots)} |1\rangle \right)$$

So if  $\omega = 0.x_1 x_2$  then we can do the following:



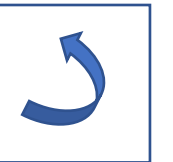
# Quantum Phase Estimation Algorithm

If  $\omega = 0.x_1x_2x_3$  then we can do the following:



# Quantum Phase Estimation Algorithm

Demo QPE





# Shor's Algorithm

## A lecture on Shor's quantum factoring algorithm

Samuel J. Lomonaco Jr.

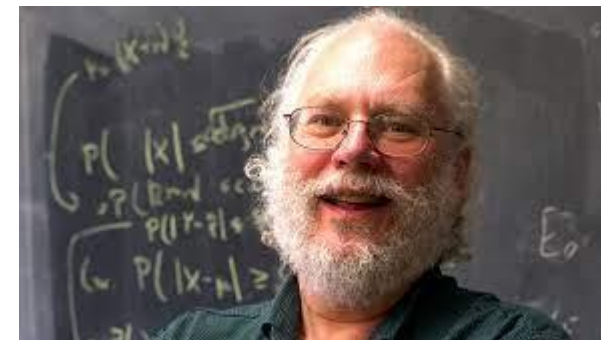
Step 1. Choose a random positive integer  $m$ . Use the polynomial time Euclidean algorithm to compute the greatest common divisor  $\gcd(m, N)$  of  $m$  and  $N$ . If the greatest common divisor  $\gcd(m, N) = 1$ , then we have found a non-trivial factor of  $N$ , and we are done. If, on the other hand,  $\gcd(m, N) \neq 1$ , then proceed to STEP 2.

Step 2. Use QUANTUM PHASE ESTIMATION to determine the unknown period  $P$  of the function

$$\begin{array}{ccc} \mathbb{N} & \xrightarrow{f_N} & \mathbb{N} \\ a & \longmapsto & m^a \bmod N \end{array}$$

Step 3. If  $P$  is an odd integer then goto Step 1. If  $P$  is even then goto Step 4.

Step 4. Since  $P$  is even, we have:  $(m^{P/2} + 1)(m^{P/2} - 1) = m^P - 1 = 0 \bmod N$   
If  $m^{P/2} + 1 = 0 \bmod N$  then goto Step 1. else goto Step 5.

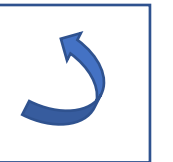


Peter Shor 1994

Step 5. Use the Euclidean algorithm to compute  $d = \gcd(m^{P/2} - 1, N)$ . Since  $m^{P/2} + 1 \neq 0 \bmod N$  it can easily be shown that  $d$  is a non-trivial factor of  $N$ . Exit with the answer  $d$ .

# Shor's Algorithm

**Demo Shor**



# The Very End!

